Electrum

Lightweight specification of behavioral models with rich configurations

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Observations

Many Alloy models feature both structural and behavioral aspects, but:

- Behavior modeling requires systematic "boilerplate"
 - explicit modeling of state (local/global state idiom)
 - every mutable construct must be indexed by state/time
 - specification of a *linear* model of time (most of the time (!))
 - specific handling of the last state of a trace
- Essentially to model check *safety* properties, indeed:
 - spurious counterexamples to *liveness* properties may happen, unless traces are enriched with *lassos* [Cunha 14, Biere *et al.* 99]
 - even then, limited to bounded model-checking (BMC)

(Safety properties rule out unwanted behaviors, liveness properties characterize expected behaviors)

Linear temporal logic (LTL):

- is more expressive than propositional logic
- is decidable
- relies on a simple & uniform model of time: infinite traces of states
- benefits from dedicated, *complete* model-checking procedures

Mark mutable fields or signatures as such (using a new var keyword).

Add LTL + primed variables (as, e.g., in TLA+).

Dedicated analyses:

- BMC by reduction to Alloy + traces with lassos
- Unbounded MC (UMC) by reduction to NuSMV or nuXmv



Example: Chord

```
sig Node {
                                       pred join [new : Node] { // an event
  var fst : lone Node.
  var snd : lone Node,
                                         new not in members
  var prdc : lone Node.
                                         some m : members {
  var todo : Status→Node }
                                            between[m, new, m.fst]
                                            fst' = fst ++ new→m.fst
var sig members in Node {}
                                            snd' = snd ++ new\rightarrowm.snd
                                            prdc' = prdc ++ new→m
var sig ringMembers in members {}
                                            todo' = todo }}
fact {
                                       fact strongFairness {
  always members =
                                         all n. m : Node {
                                            (always eventually rectifyEnabled[n.m])
    { n: Node | some n.fst and
                 some n.snd and
                                                \Rightarrow (always eventually rectify[n,m])
                                           ... }}
                 some n.prdc }
  always ringMembers =
    { m : members | m in m.^succ }} assert correctness {
                                         (eventually always not (join or fail)
fun succ : Node \rightarrow lone Node { ... }
                                           implies eventually always ideal ) }
```

Fits well most Alloy models with *behavior*.

Often leaner than plain Alloy (not always: e.g. counting events).

BMC efficiency on par with classic Alloy.

UMC with nuXmv comparable to TLA+'s TLC (room for improvement) (note: nuXmv is *not* free software; other, non-evaluated, tools exist).

Modeling [Zave 2017]'s version of Chord raised various corner cases: analyzing "abstract" liveness properties if useful (even with BMC).

- Enhance modeling of the "system" (automaton) part, *e.g.*: actions (guard + post-condition), frame rules, fairness constraints... Most models may then rely on LTL for assertions only. So add branching time (CTL) too?
- No more a conservative extension of Alloy, though.