DASH: Modelling and Analysis of Declarative State-Based Transition Systems Work in Progress (*)

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Future of Alloy Workshop Apr 2018 http://129.97.7.33:8080/dash/

(*) In International Workshop on Model-Driven Requirements Engineering (MoDRE) @ IEEE International Requirements Engineering (RE),Sep 2017 Modelling Transition Systems of Reactive Systems

Modelling the control-oriented aspects of a system can be done naturally in hierarchical and concurrent control states (*i.e.*, the Statecharts family of languages).

Control states means a state with a name.

- Control-oriented modelling languages lack abstractions for modelling data.
- We need both! An integrated model that describes both control and data aspects of the system at an abstract level.
- Goal: a natural extension to Alloy to create these integrated models.

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- Concurrent states provide a separation of concerns for aspects of the model that can happen at the "same" time.
- Events are occurences of user actions; or internal actions to model cascading effects between concurrent regions.

```
---- Allov
   abstract sig ValvePos {} <-----
1
2
   one sig OpenPos, ClosedPos extends ValvePos {}
3
   abstract sig Room {}
4
   5
       valve: Room -> ValvePos <----- snapshot variables
6
7
       desiredTemp: Room -> Int
8
       . . .
9
       event deactivate {} _____
       init { <----- initial snapshot constraint</pre>
10
           all r: Room | r.valve = ClosedPos <---- Allov
11
12
       }
13
       conc state Furnace { ...
                                     }
14
       conc state Controller {
15
           default state Off { }
           state On { <----- control state</pre>
16
17
               trans t3 { <---- transition
18
                   on heatSwitchOff
                                     <----- event
19
                   goto Off
                   do r.valve' = ClosedPos <-- Alloy (w/ prime)
20
21
                   send deactivate
                                    <---- send an event
22
               }
23
            . . .
                                       HeatingSystem
24
       }
                                       Furnace Controller
25
        . . .
                                              Off
   }
26
                                               t3
                                              On
http://129.97.7.33:8080/dash/
                                                           7
```

Syntax of Core Dash

Everything in Alloy plus:

```
[conc|default] state <name> { ... }
1
2
3
   trans <name> {
4
        from <src_state>
        on <trigger_event> <--- on/when = precondition
5
6
        when <alloy>
7
        goto <dest_state>
        do <alloy> <----- postcondition
8
9
        send <generated_event>
10
   }
11
12
    event <name> {}
13
14
    condition <name> {}
```

Transitions understood within context (source state, etc.)

- Concise
- Permits factoring by state, event, and condition
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Add-ons:
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addon (do incErrorCounter)
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Transition templates (reusuable)

Semantics of Dash

- A "big" step (model's response to env) can be multiple transitions in different concurrent states.
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- A "big" step (model's response to env) can be multiple transitions in different concurrent states.
- Philosophy for Semantic Decisions: Atomicity of transitions; internal sequencing of transitions in concurrent states should be rare.
- Using the framework of Esmaeilsabzali:
 - Concurrency: only one concurrent state can take a transition in a big step (avoids race conditions and makes model's behaviour more understandable).
 - Big-step Maximality: only one transition per concurrent region in a big step (avoids potentially infinite big-steps).
 - **Event Lifeline:** events last the entire big step.
 - Variable Lifeline: variables change their values in small steps.
 - Priority: Outer state over inner state.

Semantics: Frame Problem

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Our solution:

- If variable a is NOT designated environmental and is not mentioned in the action of the taken transition, it keeps its value from the previous snapshot.
- ► To override this semantics, action can be that variable a must be within its range of values (0 ≤ a' ≤ 100).

Tool Chain: xText

- Grammar for Dash (which includes Alloy)
- "Smart" editor automatically created
- Transformation rules to:
 - 1. Transform to Core Dash (state hierarchy plus fully detailed transitions)
 - 2. Transform Core Dash to Alloy
 - Transition names are present in Alloy so it's easier to relate a counterexample to the model.

Analyzing Transition Systems in Alloy

BMC (Bounded Model Checking) Scoped TCMC (Transitive-Closure-based Model Checking for CTLFC) Significance axioms to ensure the TS is "big enough"

Table: Deducing Complete Model Checking Results

		Scoped TCMC		BMC using ordering	
Property		Pass	Fail	Pass	Fail
Safety		Ambiguous	Real Bug	Ambiguous	Real Bug
Finite	w/o dead-loop	Ambiguous	Real Bug	Real Pass	Ambiguous
Liveness	w/ dead-loop	Real Pass	Ambiguous		
Infinite Liveness		Ambiguous	Real Bug	Cannot Express	
Existential		Real Pass	Ambiguous	Cannot Express	

Summary

- Dash is an extension of Alloy to create an integrated model that describes both control and data aspects of the system at an abstract level.
- Dash provides explicit syntax for:
 - Transitions, preconditions, postconditions
 - Hierarchical and concurrent control states
 - Events to allows cascading effect between concurrent regions
 - Semantics chosen to match combination of declarative and control-oriented modelling paradigms and address the frame problem.
 - Syntactic sugar for conciseness (transition comprehension, add-ons, factoring by events, conditions, transition templates).
- Dash is fully integrated with Alloy for everything else (expressions, data abstractions, functions, and predicates).